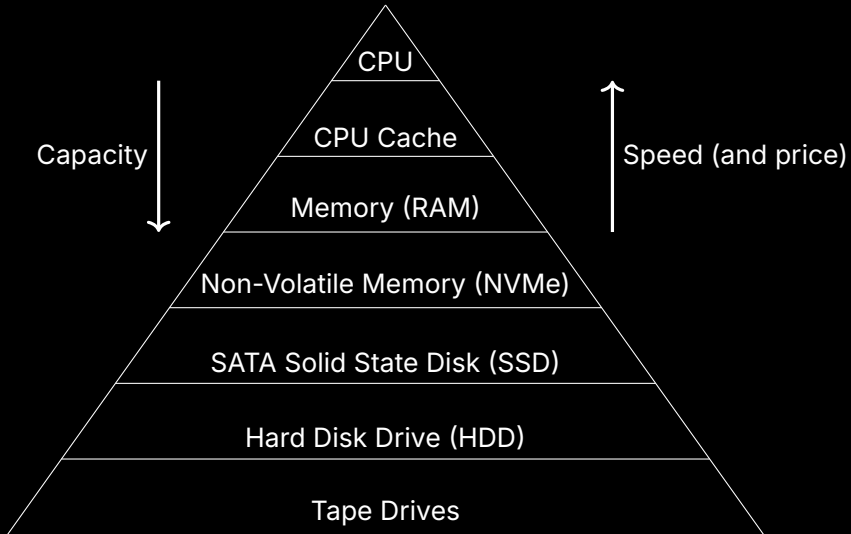


Page Replacement

2024 Fall ECE 344: Operating Systems
Jon Eyolfson

Lecture 29
2.0.1

Computer Memory Hierarchy is a Trade-off of Capacity and Speed



We Want to Hide the Hierarchy from the User

Each level wants to pretend it has the speed of the layer above it
and the capacity of the layer below it

The memory used by all the processes may exceed the amount of physical
memory

Not all of them may be in use at the same time

Only keep referenced pages in memory, put others on disk
Swap pages back to memory when they're needed

Demand Paging

We use memory as a cache for the file system

Map memory pages to file system blocks

Only load them into memory when they're used through the page fault handler

If the page doesn't represent a file, we can map it to swap space

Move it to disk if we need to use more physical memory

A Processes' Working Set Should Fit in Physical Memory

Given an amount of time, the number of pages your process uses is called its *working set*

If you cannot fit your working set into physical memory your process will *thrash*

It's constantly moving entries in and out of cache

Page Replacement Algorithms

Optimal

Replace the page that won't be used for the longest

Random

Replace a random page

First-in First-out (FIFO)

Replace the oldest page first

Least Recently Used (LRU)

Replace the page that hasn't been used in the longest time

Page Replacement Evaluation

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

We'll use this for a bunch of examples during this lecture

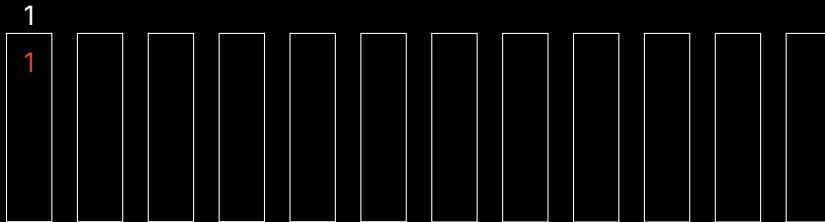
We want the fewest number of page faults

For every example we'll find the number of page faults

Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

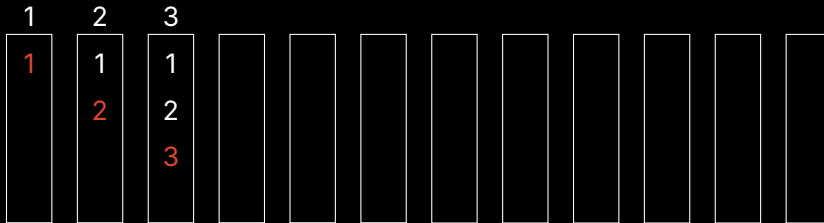
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

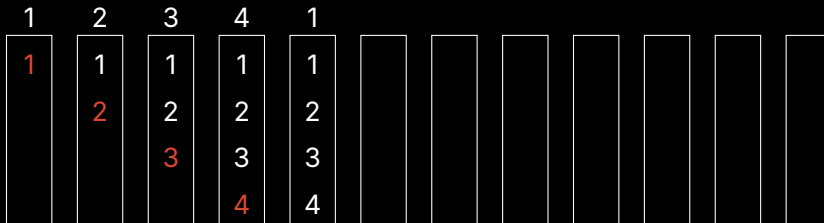
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

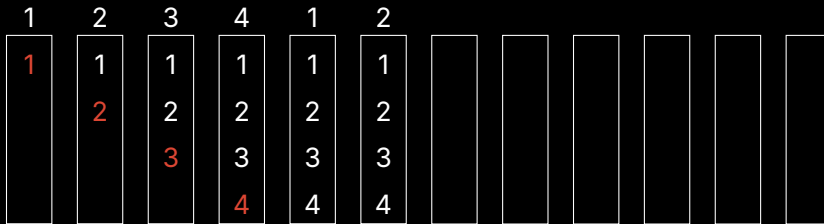
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

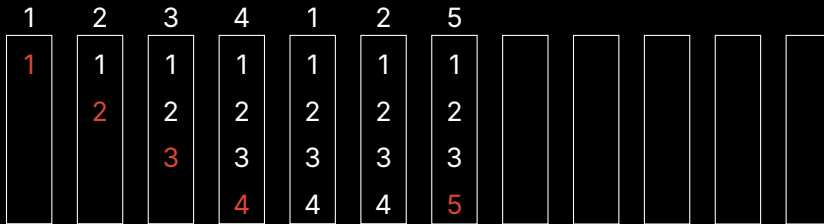
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

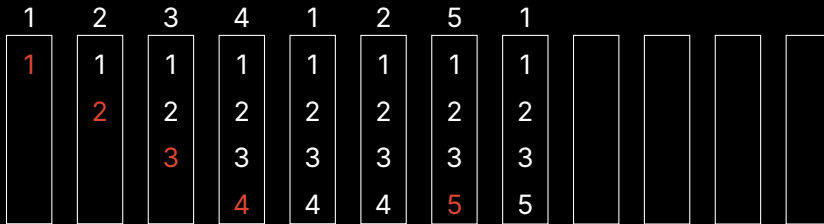
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

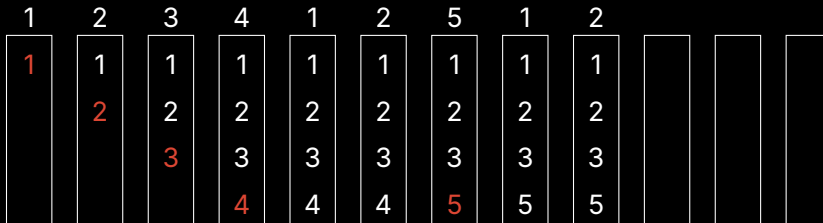
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

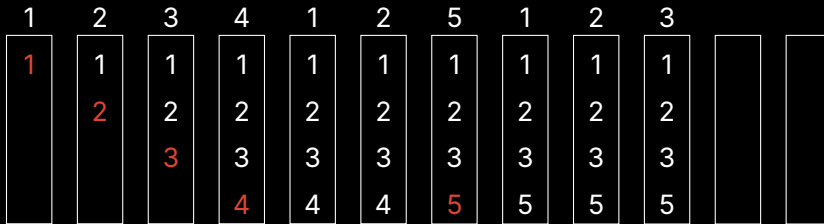
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

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Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	
1	1	1	1	1	1	1	1	1	1	4	
	2	2	2	2	2	2	2	2	2	2	
		3	3	3	3	3	3	3	3	3	
			4	4	4	5	5	5	5	5	

Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	5
1	1	1	1	1	1	1	1	1	1	4	4
	2	2	2	2	2	2	2	2	2	2	2
		3	3	3	3	3	3	3	3	3	3
			4	4	4	5	5	5	5	5	5

Optimal Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

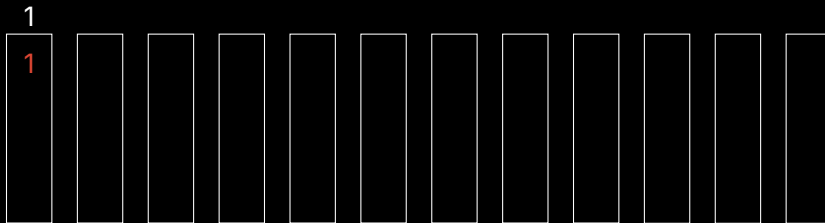
1	2	3	4	1	2	5	1	2	3	4	5
1	1	1	1	1	1	1	1	1	1	4	4
	2	2	2	2	2	2	2	2	2	2	2
		3	3	3	3	3	3	3	3	3	3
			4	4	4	5	5	5	5	5	5

6 page faults

FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

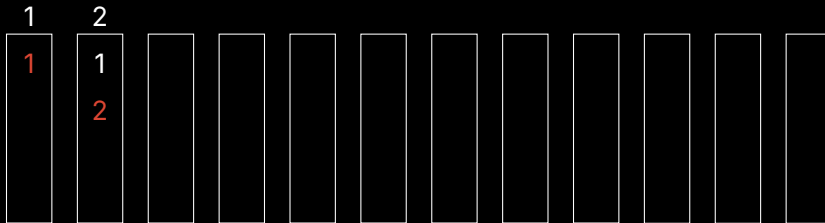
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

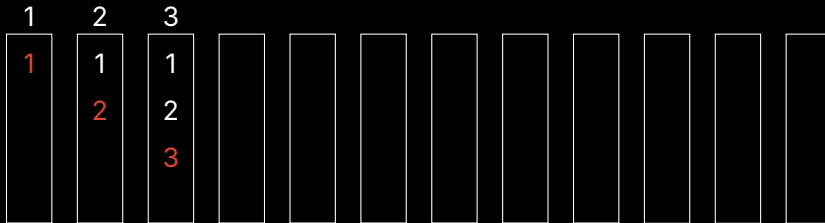
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

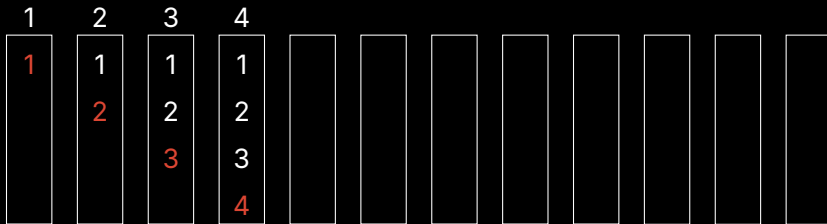
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

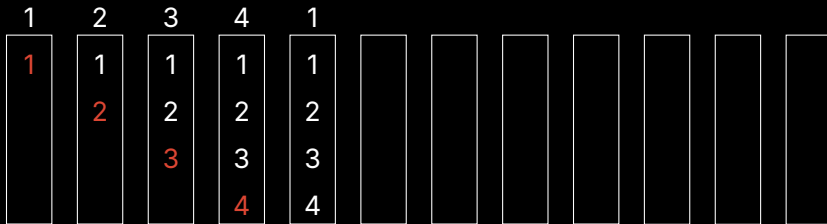
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

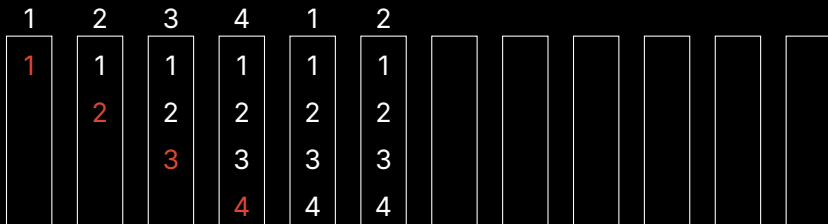
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

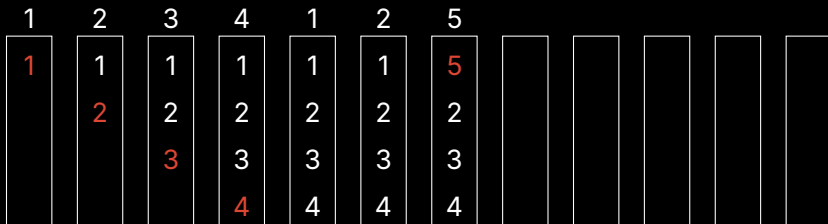
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

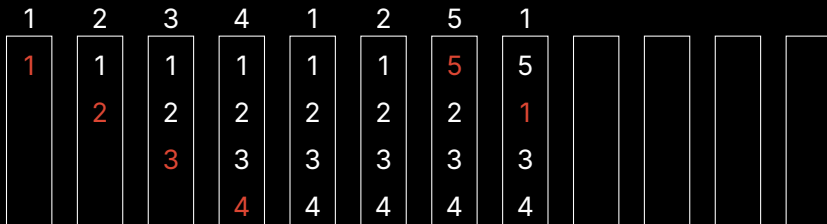
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

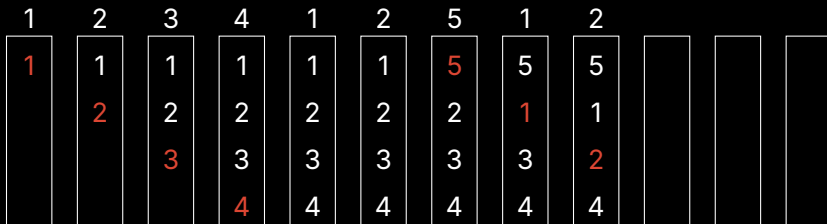
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

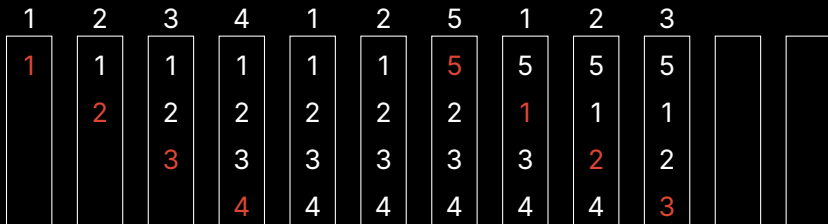
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	
1	1	1	1	1	1	5	5	5	5	4	
	2	2	2	2	2	2	1	1	1	1	
		3	3	3	3	3	3	2	2	2	
			4	4	4	4	4	4	3	3	

FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	5
1	1	1	1	1	1	5	5	5	5	4	4
	2	2	2	2	2	2	1	1	1	1	5
		3	3	3	3	3	3	2	2	2	2
			4	4	4	4	4	4	3	3	3

FIFO Example

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	5
1	1	1	1	1	1	5	5	5	5	4	4
	2	2	2	2	2	2	1	1	1	1	5
		3	3	3	3	3	3	2	2	2	2
			4	4	4	4	4	4	3	3	3

10 page faults

FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

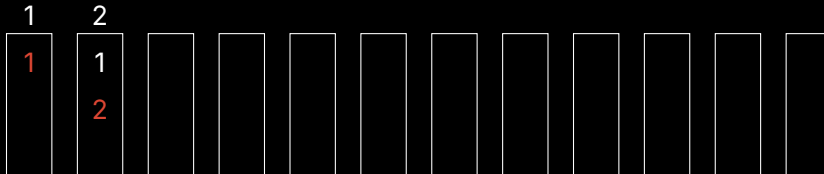
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



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1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



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1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	5
1	1	1	4	4	4	5	5	5	5	5	5
	2	2	2	1	1	1	1	1	3	3	3
		3	3	3	2	2	2	2	2	4	4

FIFO Example (3 Page Frames)

Assume our physical memory can only hold **3** pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	5
1	1	1	4	4	4	5	5	5	5	5	5
	2	2	2	1	1	1	1	1	3	3	3
		3	3	3	2	2	2	2	2	4	4

9 page faults

Bélády's Anomaly Says More Page Frames Causes More Faults

This is a problem with FIFO algorithms

Does not exist with LRU or "stack-based algorithms"

Paper in 2010 found that this FIFO anomaly is unbounded

(<https://arxiv.org/abs/1003.1336>)

You could construct a sequence to get any arbitrary page fault ratio

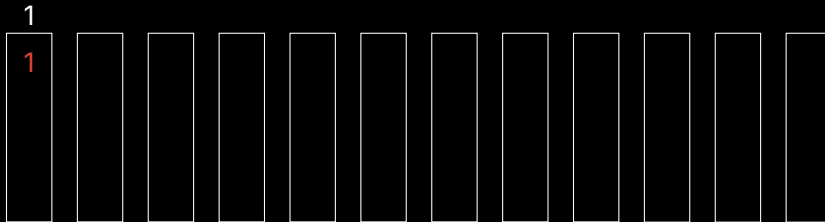
For other algorithms:

increasing the number of page frames decreases the number of page faults

LRU Example (Use FIFO to Break Ties)

Assume our physical memory can only hold 4 pages, and we access the following:

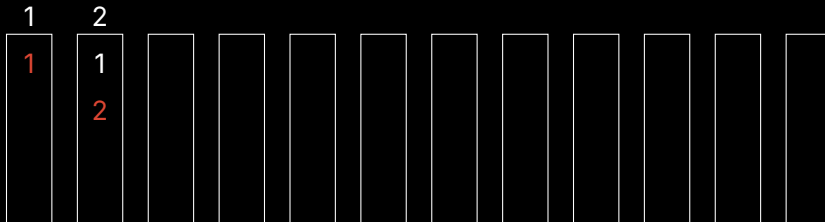
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



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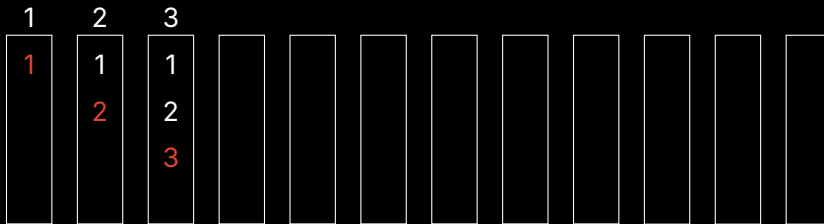
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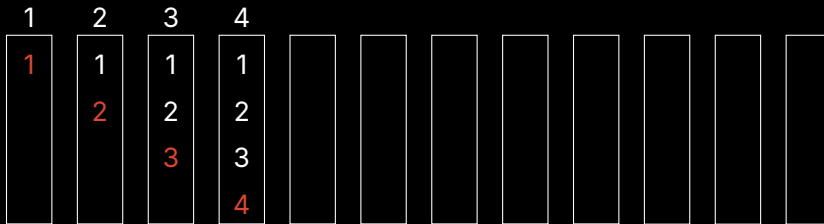
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



LRU Example (Use FIFO to Break Ties)

Assume our physical memory can only hold 4 pages, and we access the following:

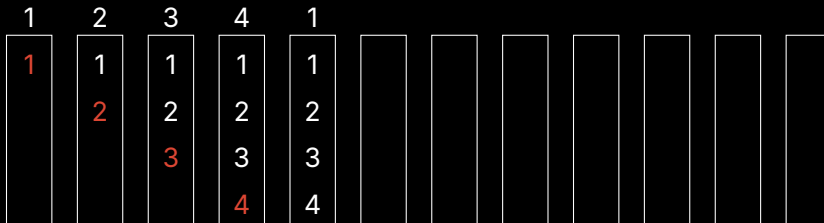
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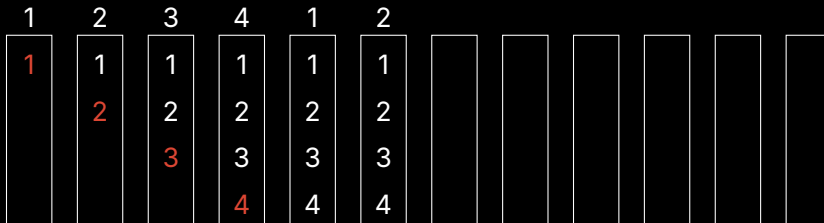
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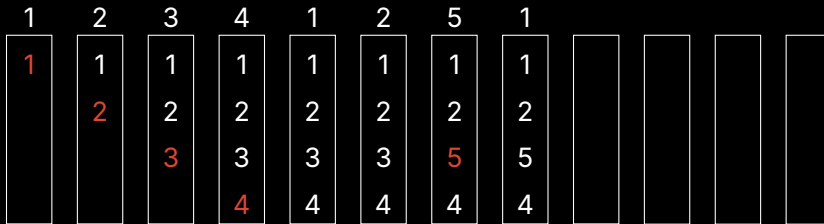
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



LRU Example (Use FIFO to Break Ties)

Assume our physical memory can only hold 4 pages, and we access the following:

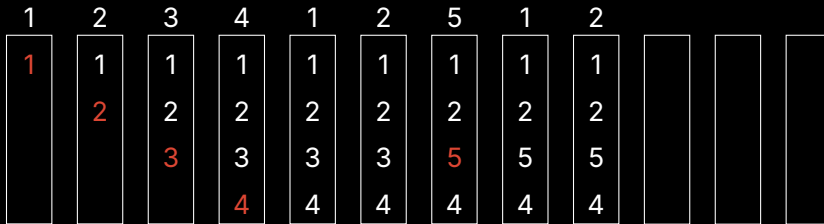
1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)



LRU Example (Use FIFO to Break Ties)

Assume our physical memory can only hold 4 pages, and we access the following:

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LRU Example (Use FIFO to Break Ties)

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	
1	1	1	1	1	1	1	1	1	1	1	
	2	2	2	2	2	2	2	2	2	2	
		3	3	3	3	5	5	5	5	4	
			4	4	4	4	4	4	3	3	

LRU Example (Use FIFO to Break Ties)

Assume our physical memory can only hold 4 pages, and we access the following:

1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

1	2	3	4	1	2	5	1	2	3	4	5
1	1	1	1	1	1	1	1	1	1	1	5
	2	2	2	2	2	2	2	2	2	2	2
		3	3	3	3	5	5	5	5	4	4
			4	4	4	4	4	4	3	3	3

LRU Example (Use FIFO to Break Ties)

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1 2 3 4 1 2 5 1 2 3 4 5 (all of the pages are initially on disk)

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1	1	1	1	1	1	1	1	1	1	1	5
	2	2	2	2	2	2	2	2	2	2	2
		3	3	3	3	5	5	5	5	4	4
			4	4	4	4	4	4	3	3	3

8 page faults

Implementing LRU in Hardware Has to Search All Pages

You could implement it by keeping a counter for each page

For each page reference, save the system clock into the counter

For replacement, scan through the pages and find the one with the oldest clock

Implementing LRU in Software is Too Expensive

Create a doubly linked list of pages

For each page reference, move it to the front of the list

For replacement, remove from the back of the list

It requires 6 pointer updates for each page reference, and
also creates a high contention bottleneck for multiple processors

Implementing LRU in Practice Isn't Going to Work

We settle for approximate LRU

- LRU is an approximation of the optimal case anyways

There's lots of different tweaks you can do to implement it more efficiently

We'll be looking at the clock algorithm, but there's also:

- Least Frequently Used (LFU), 2Q, Adaptive Replacement Cache (ARC)

Page Replacement Algorithms Aim to Reduce Page Faults

We saw the following:

- Optimal (good for comparison but not realistic)
- Random (actually works surprisingly well, avoids the worst case)
- FIFO (easy to implement but Bélády's anomaly)
- LRU (gets close to optimal but expensive to implement)