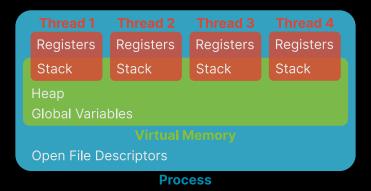
User Threads Lab Primer and Sockets

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Lecture 12

A Process with 4 Threads



Remember, You Can Access These Examples!

Look at the materials repository, today: lectures/12-user-threads-lab-primer-and-sockets

First Example

ucontext-example.c

Done live, please see the recording

Second Example

tailq-example.c

Done live, please see the recording

ucontext Question

```
Global variables: int i = 0; ucontext_t uA; // Initialized to execute thread_a() ucontext_t uB; // Initialized to execute thread b()
```

One kernel thread calls set_context(&uA), what happens?

```
void thread_a() {
                                         void thread_b() {
    int d = 0;
                                              int d = 1;
    while (i < 3) {
                                              while (i < 3) {
        i++;
                                                  i++;
        printf("A: %d\n", i);
                                                  printf("B: %d\n", i);
        d = 0;
                                                  d = 1;
        getcontext(&uA);
                                                  getcontext(&uB);
        if (d = 0) {
                                                  if (d = 1) {
                                                      d = 0:
            d = 1:
            setcontext(&uB);
                                                      setcontext(&uA);
```

Sockets are Another Form of IPC

We've seen pipes and signals
We also talked about shared memory

Previous IPC assume that the processes are on the same physical machine

Sockets enable IPC between physical machines, typically over the network

Servers Follow 4 Steps to Use Sockets

These are all system calls, and have the usual C wrappers:

- 1. socket
 - Create the socket
- 2. bind
 - Attach the socket to some location (a file, IP:port, etc.)
- 3. listen Indicate you're accepting connections, and set the queue limit
- 4. accept
 - Return the next incoming connection for you to handle

Clients Follow 2 Steps to Use Sockets

Clients have a much easier time, they use one socket per connection

- 1. socket
 - Create the socket
- 2. connect

Connect to some location, the socket can now send/receive data

The socket System Call Sets the Protocol and Type of Socket

```
int socket(int domain, int type, int protocol);
domain is the general protocol, further specified with protocol (mostly unused)
    AF_UNIX is for local communication (on the same physical machine)
    AF_INET is for IPv4 protocol using your network interface
    AF_INET6 is for IPv6 protocol using your network interface
type is (usually) one of two options: stream or datagram sockets
```

Stream Sockets Use TCP

All data sent by a client appears in the same order on the server Forms a persistent connection between client and server Reliable, but may be slow

Datagram Sockets Use UDP

Sends messages between the client and server No persistent connection between client and server Fast but messages may be reordered, or dropped

The bind System Call Sets a Socket to an Address

socket is the file descriptor returned from the socket system call

```
There's different sockaddr structures for different protocols struct sockaddr_un for local communcation (just a path) struct sockaddr_in for IPv4, a IPv4 address (e.g. 8.8.8) struct sockaddr_in6 for IPv6, a IPv6 address (e.g. 2001:4860:4860:8888)
```

The listen System Call Sets Queue Limits for Incoming Connections

int listen(int socket, int backlog);
socket is still the file descriptor returned from the socket system call
backlog is the limit of the outstanding (not accepted) connections
 The kernel manages this queue, and if full will not allow new connections
We'll set this to 0 to use the default kernel queue size

The accept System Call Blocks Until There's a Connection

socket is *still* the file descriptor returned from the socket system call address and address_len are locations to write the connecting address. Acts as an optional return value, set both to NULL to ignore. This returns a new file descriptor, we can read or write to as usual

The connect System Call Allows a Client to Connect to an Address

sockfd is the file descriptor returned by the socket system call
The client would need to be using the same protocol and type as the server
addr and addrlen is the address to connect to, exactly like bind
If this call succeeds then sockfd is may be used as a normal file descriptor

Our Example Server Sends "Hello there!" to Every Client and Disconnects

Please see lectures/19-sockets in your materials repository Relevant source files: client.c and server.c

We use a local socket just for demonstration, but you could use IPv4 or IPv6 We use example.sock in the current directory as our socket address

Our server uses signals to clean up and terminate from our infinite accept loop

Instead of read/write There's Also send/recv System Calls

These system calls are basically the same thing, except they have flags

Some examples are:

MSG_00B — Send/receive out-of-band data

MSG_PEEK — Look at data without reading

MSG_DONTROUTE — Send data without routing packets

Except for maybe MSG_PEEK, you do not need to know these

sendto/recvfrom take an additional address

The kernel ignores the address for stream sockets (there's a connection)

You Perform Networking Through Sockets

Sockets are IPC across physical machines, the basics are:

- Sockets require an address (e.g. local and IPv4/IPv6)
- There are two types of sockets: stream and datagram
- Servers need to bind to an address, listen, and accept connections
- Clients need to connect to an address